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# KAPREKAR NUMBERS AND ITS ANALOG EQUATIONS 

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#### Abstract

The present article discusses the contribution of D.R.Kaprekarin a recreational number theory, particularly Kaprekar Constant and Kaprekar Number. The analog equations of Kaprekar number 1. $K^{3}=P_{0}+P_{1}(10)+P_{2}\left(10^{2}\right)+P_{3}\left(10^{3}\right)+\ldots+P_{n}\left(10^{n}\right)$ $\mathrm{K}=P_{0}+P_{1}+P_{2}+P_{3}+\ldots+P_{n}$ 2. $K^{3}=P\left(10^{n}\right)-Q ; 0 \leq Q \leq P$ $\mathrm{K}=P-Q$ 3. $\mathrm{K}^{3}=P\left(10^{n}\right)+Q$ $\mathrm{K}=\mathrm{P}+\mathrm{Q}$; where P and Q are positive integers, $0<\mathrm{Q}<\mathrm{P}$ and $\mathrm{n} \geq 1$. are discussed. C Programs are developed for the same and for Kaprekar Number the program is dynamic is enough to suit any base.


Keywords: Kaprekar Constant, Kaprekar Number, Kaprekar Cycle, Reverse subtraction process.

## 1. INTRODUCTION

D.R.Kaprekar was born on January 17, 1905 in Maharashtra. For whatever reason he is not so well known in Mathematical world but his contribution in recreation mathematics cannot be ignored. He discovered Kaprekar Constant, Kaprekar number, Self-numbers, Harshad numbers, Demlo numbers and many more such numbers. In this paper C program for Kaprekar Constant, Kaprekar number and its analog equations are discussed which can be considered as a tribute to D.R.Kaprekar. The C program developed in this paper is dynamic enough to get Kaprekar number for any base 10, 2(Binary), 8(Octal) and so on.

## 2. KAPREKAR CONSTANT AND KAPREKAR PROCESS

## Reverse Subtraction Process

If we consider any 4-digit number where all the digits should not be alike, then we can generate two different numbers by arranging them into ascending and descending order. Subtraction of smaller number from larger number will generate another four digit number. Occasionally, a three digitsnumber will be generated then put zero on the extreme left side of the number. This process is called as reverse subtraction process.

Now the question arises that if we keep on repeating the above said reverse subtraction process, Can we have an infinite sequence of different numbers? Or a sequence of finite numbers.

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Let us consider the following example:

## Number = 3947

Step 1: Number = 3947
Ascending of Number: 3479
Descending of Number: 9743
Reverse Subtraction Process: 6264
Step 2: Number = 6264
Ascending of Number: 2466
Descending of Number: 6642
Reverse Subtraction Process: 4176
Step 3: Number $=4176$
Ascending of Number: 1467
Descending of Number: 7641
Reverse Subtraction Process: 6174
Step 4: Number = 6174
Ascending of Number: 1467
Descending of Number: 7641
Reverse Subtraction Process: 6174
Now, there is no use of repeating the process as step3\& step4 are same. Therefore starting from 3947, a finite sequence of numbers \{3947, 6264, 4176and 6174\} will be generated with the help of Reverse Subtraction Process. The interesting fact is that if we start with any four digit number (not all the digits are same) then within eight or less steps the sequence converges to 6174, which is known as Kaprekar Constant. The same is developed here in C program.

## Program for Kaprekar Process [Kaprekar Constant]

```
#include <stdio.h>
#include <conio.h>
#include <alloc.h>
#include <values.h>
voidAscendingOrder(int ascending[],int n)
{
    inti,k,y;
    for(k=1;k<n;k++)
    {
        y=ascending[k];
        for(i=k-1;i>=0 && y<ascending[i]; i--)
            ascending[i+1]=ascending[i];
        ascending[i+1]=y;
    }
}
voidDescendingOrder(int descending[],int n)
{
    inti,k,y;
    for(k=1;k<n;k++)
    {
        y=descending[k];
        for(i=k-1;i>=0 && y>descending[i]; i--)
                        descending[i+1]=descending[i];
        descending[i+1]=y;
    }
}
//Function is used to conver number into array
int * numberToArray(int number)
{
    inti = 0,*arr=NULL;
    arr=(int *)malloc(4);
    if(number<1000 && number>99)
    {
```

```
arr[i]=0;
        i++;
    }
    if(number<100 && number>9)
    {
        arr[i]=0;
        arr[++i]=0;
        i++;
    }
    if (number>1 && number<10)
    {
        arr[i]=0;
        arr[++i]=0;
        // i++;
    }
    while (number > 0)
    {
        arr[i] = number % 10;
        number /= 10;
        i++;
    }
    returnarr;
}
void main(void)
{
    inti,sum=0,cnt=0,temp=MAXINT;
    int *arr,*arr1,ans1[10],n,ans,n1=0,n2=0;
    clrscr();
    printf("Enter a Number:");
    scanf("%d",&n);
    arr = numberToArray(n);
con:
    arr1 = arr;
    printf("\n");
    DescendingOrder(arr,4);
    for(i=0;i<4;i++)
    {
        printf(" %d ",arr[i]);
        n1 = 10 * n1 + arr[i];
}
printf("\n");
AscendingOrder(arr1,4);
for(i=0;i<4;i++)
{
    printf(" %d ",arr1[i]);
    n2 = 10 * n2 + arr1[i];
}
printf("\n");
ans = n1-n2;
printf("\n Answer is %d",ans);
getch();
if (ans<1000)
{ }
if (ans==temp)
gotoexi;
else
{
cnt++;
temp=ans;
*arr=NULL;
n1=n2=0;
arr = numberToArray(ans);
    goto con;
}
```

```
exi:
printf("\n No. of Iterations %d",cnt);
        getch();
}
/* OUTPUT
Enter a Number: }123
4 21
1234
Answer is 3087
8 3 0
0 3 8
Answer is 8352
8 3 2
2 35 8
Answer is 6174
7641
1467
Answer is 6174
No. of Iterations 3
*
/
3. KAPREKAR CYCLE
```

Note that the above Kaprekar process is not applicable for single digit number but it is applicable for more than one digit numbers. For two digit number the Kaprekar Process is:

## Number: 89

1. Ascending Number : 98 Descending Number : 89
Subtraction : 09
2. Ascending Number : 09

Descending Number : 90
Subtraction : 81
3. Ascending Number : 18

Descending Number : 81
Subtraction : 63
4. Ascending Number : 36

Descending Number: 63
Subtraction : 27
5. Ascending Number : 27

Descending Number: 72
Subtraction : 45
6. Ascending Number : 45

Descending Number : 54
Subtraction : 09
Note that step 2 \& step 6 are same. This process generates Kaprekar Cycle $89 \rightarrow 09 \rightarrow 81 \rightarrow 63 \rightarrow 27 \rightarrow 45 \rightarrow 09$. Similarly, one can have a Kaprekar Cycle for other digits also but for four digits we get Kaprekar cycle of length 1. Amazing fact is Kaprekar constant or number belonging to Kaprekar Cycle are divisible by 9, see [5].

## 4. KAPREKAR NUMBER

Number $K$ is said to be Kaprekar number if $K^{2}$ is divided into two parts, left and right, such that the sum of the two parts is equal to $K$. Mathematically, $K$ is a Kaprekar number if

$$
K^{2}=P * 10^{n}+Q\left(n>=1, P>=0.0<Q<10^{n}\right)
$$

and

$$
K=P+Q
$$

Therefore,

$$
\begin{aligned}
& (P+Q)^{2}=P\left(10^{n}\right)+Q \\
\Rightarrow & P^{2}+2 P Q+Q^{2}-\left(10^{n}\right) P+Q=0 \\
\Rightarrow & P^{2}+P\left(2 Q-10^{n}\right)+Q^{2}+Q=0 \\
\Rightarrow & P=\frac{10^{n}-2 Q \pm \sqrt{10^{2 n}-4 Q\left(10^{n}\right)+4 Q}}{2}
\end{aligned}
$$

Also,

$$
\begin{aligned}
& 10^{2 n}-4 Q\left(10^{n}-1\right) \geq 0 \\
\Rightarrow & 10^{2 n} \geq 4 Q\left(10^{n}-1\right) \\
\Rightarrow & 0 \leq Q \leq \frac{10^{2 n}}{4\left(10^{n}-1\right)}, \text { give the upper bound for } Q .
\end{aligned}
$$

For the fixed value of $n$, the $C$ program developed is based on the following

1. $10^{2 n}-4 Q\left(10^{n}-1\right)$ is a perfect square
2. The upper bound of Q is $\frac{10^{2 n}}{4\left(10^{n}-1\right)}$.

The size of data type in C depends on machine so the existing program execution depends on the machine. Some examples of Kaprekar numbers are
$9^{2}=81$ and $9=8+1$
$5050^{2}=25502500$ and $5050=2550+2500$
$703^{2}=494209$ and $703=494+209$
Observed that for any Kaprekar number there exists another Kaprekar number such that $K_{1}+K_{2}=10^{n}$ where n is a positive integer.

## Program for Kaprekar Numbers

\#include<stdio.h>
\#include<conio.h>
\#include<math.h>
\#include<process.h>
void main()
\{

```
double n,q,div,div1,squareroot,p,p1,perfectsqrt[100],qq[100];
inti=0,j=0,k,k1,cnt=0,iteration,psqrt;
clrscr();
printf("Enter the value of n:");
scanf("%lf",&n);
div = pow(10,2*n);
div1 = 4*(pow(10,n)-1);
iteration = div / div1;
    for(q=0;q<iteration;q++)
    {
                squareroot = sqrt(div-(q*div1));
                psqrt = squareroot;
                if(squareroot == psqrt)
                {
                qq[i++] = q;
                perfectsqrt[j++] = squareroot;
                        cnt++;
}
```

iteration ++;

```
    }
for (i=0;i<cnt;i++)
{
    p=(pow(10,n)-(2*qq[i])+perfectsqrt[i])/2;
    p1=(pow(10,n)-(2*qq[i])-perfectsqrt[i])/2;
    k=p + qq[i];
    k1 = p1 + qq[i];
    printf("\nPerfect Squares (P + Q) is %d = %lf + %0.2lf",k,p,qq[i]);
    printf("\nPerfect Squares (P + Q) is %d = %lf + %0.2lf",k1,p1,qq[i]);
}
getche();
}
/*OUTPUT
Enter the value of n : 3
Perfect Squares \((\mathrm{P}+\mathrm{Q})\) is \(1000=1000.000000+0.00\)
Perfect Squares \((\mathrm{P}+\mathrm{Q})\) is \(0=0.000000+0.00\)
Perfect Squares \((\mathrm{P}+\mathrm{Q})\) is \(999=998.000000+1.00\)
Perfect Squares \((\mathrm{P}+\mathrm{Q})\) is \(1=0.000000+1.00\)
Perfect Squares \((\mathrm{P}+\mathrm{Q})\) is \(703=494.000000+209.00\)
Perfect Squares \((\mathrm{P}+\mathrm{Q})\) is \(297=88.000000+209.00\)
*/
```

The above program can be used for any other base by just replacing 10 by the required base number. The following are the sample outputs for the binary base 2 .

## Sample Output for binary base

Enter the value of n : 3
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $8=8.000000+0.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $0=0.000000+0.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $7=6.000000+1.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $1=0.000000+1.00$
Enter the value of n : 4
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $16=16.000000+0.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $0=0.000000+0.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $15=14.000000+1.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $1=0.000000+1.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $10=6.000000+4.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $6=2.000000+4.00$
Enter the value of n : 10
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $1024=1024.000000+0.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $0=0.000000+0.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $1023=1022.000000+1.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $1=0.000000+1.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $837=684.000000+153.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $187=34.000000+153.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $682=454.000000+228.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $342=114.000000+228.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $528=272.000000+256.00$
Perfect Squares $(\mathrm{P}+\mathrm{Q})$ is $496=240.000000+256.00$
*/
Some of the examples for binary base are

1. $7=(6)_{2}+1_{2}$ and $7^{2}=49=(6)_{2}\left(2^{3}\right)+1_{2}$ that is $7=110+001$ and $7^{2}=49=110001$
2. $837=684_{2}+153_{2}$ and $837^{2}=700569=(684)_{2}\left(2^{10}\right)+153$ $837=1010101100+10011001$ and $837^{2}=101010110010011001$

Observed that for any Kaprekar number with base 2 there exists another Kaprekar number with base 2 such that $K_{1}+K_{2}=2^{n}$, where $n$ is a positive integer.

## 5. ANALOG EQUATIONS OF KAPREKAR NUMBER

$5.1 K^{3}=P_{0}+P_{1}(10)+P_{2}\left(10^{2}\right)+P_{3}\left(10^{3}\right)+\ldots+P_{n}\left(10^{n}\right)$
$K=P_{0}+P_{1}+P_{2}+P_{3}+\ldots+P_{n}$

## Program for 5.1

\#include<stdio.h>
\#include<conio.h>
\#include<math.h>
void main()
\{
long double x[150];
longinti, sum,n,a[150]; clrscr();
printf("Enter the value of n");
scanf("\%ld",\&n);
for( $\mathrm{i}=0 ; \mathrm{i}<=\mathrm{n} ; \mathrm{i}++$ )
\{ $a[i]=p o w(i, 3)$;
while(a[i]>0)
\{
$x[i]=\operatorname{fmod}(a[i], 10)$;
sum $=$ sum $+\mathrm{x}[\mathrm{i}]$;
$\mathrm{a}[\mathrm{i}]=\mathrm{a}[\mathrm{i}] / 10$; \}
if(sum == i) printf("ln i = \%ld sum = \%ld",i,sum);
sum $=0$;
\} getche();
\}
/*OUTPUT
Enter the value of n100
$\mathrm{i}=1$ sum $=1$
$\mathrm{i}=8$ sum $=8$
$\mathrm{i}=17$ sum $=17$
$\mathrm{i}=18$ sum $=18$
$\mathrm{i}=26$ sum $=26$
$\mathrm{i}=27$ sum $=27$
*/
Observed that
$8^{3}=512$ and $8=5+1+2$
$26^{3}=17576$ and $26=1+7+5+7+6$
Surprisingly, such numbers are very few and above output gives all such numbers.
$5.2 K^{3}=P\left(10^{n}\right)-Q ; 0 \leq Q \leq P$
$K=P-Q$
Therefore,

$$
\begin{aligned}
& K^{3}-K=P\left(10^{n}-1\right) \\
\Rightarrow & K(K-1)(K+1)=P\left(10^{n}-1\right)
\end{aligned}
$$

Thus, all the Kaprekar numbers will satisfy the above mentioned equation. Some of the examples are:
$55^{3}=166375=1680\left(10^{2}\right)-1625$ and $55=1680-1625$
$703^{3}=347428927=347776\left(10^{3}\right)-347073$ and $703=347776-347073$
Clearly, the density of such numbers may be more than Kaprekar Numbers since there may exists numbers which are not Kaprekar number but may satisfy the above equation.
5.3 $K^{3}=P\left(10^{n}\right)+Q$
$K=P+Q$; where, $P$ and $Q$ are positive integers, $0<Q<P$ and $n \geq 1$.
Therefore, $(P+\mathrm{Q})^{3}=P\left(10^{n}\right)+Q$

$$
\begin{aligned}
& P^{3}+Q^{3}+3 P^{2} Q+3 P Q^{2}-P\left(10^{n}\right)-Q=0 \\
& P^{3}+3 P^{2} Q+P\left[3 Q^{2}-10^{n}\right]+Q^{3}-Q=0
\end{aligned}
$$

Sum of roots $=-3 Q$
Diminishing the roots by $-Q$


Therefore, transformed equation is

$$
y^{3}-10^{n} y+Q\left(10^{n}-1\right)=0
$$

Let $y=u+v$
Therefore, $y^{3}-3 u v y-\left(u^{3}+v^{3}\right)=0$
Comparing, $u^{3}+v^{3}=10^{3 n} / 27$ and $u^{3}+v^{3}=-Q\left(10^{n}-1\right)$
If $\mathrm{u}^{3}$ and $v^{3}$ are the root of equation $t^{2}+Q\left(10^{n}-1\right)+\frac{10^{3 n}}{27}=0$
$\Rightarrow t=\frac{-Q\left(10^{n}-1\right) \pm \sqrt{Q^{2}\left(10^{n}-1\right)^{2}-4 \frac{10^{3 n}}{27}}}{2}$
Here, $10^{3 \mathrm{n}} / 27$ cannot be an integer, hence, no positive integers P and Q exists for which equation 5.3 is satisfied except $\mathrm{Q}=0$.

## 6. CONCLUDING REMARKS

In the Kaprekar process, the most amazing part is the number 6174. But if one thinks logically then for any n digit number if one can able to define the process(Reverse Addition, division and so on.) which generates another n digit number uniquely then process has to terminate after finitely many steps, since we have finite $n$ digit numbers. It means the sequence has to hit a previously generated number. In any case we can have a cycle or a constant. Coincidently, for four digit number the reverse subtraction process gives a cycle of length 1.The analog equations for the Kaprekar numbers are defined and discussed. Many more such equations can be defined in future.

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